# MINIX 3 Kernel API

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## Abstract

In general, kernel calls allow system processes to request kernel services, for example, to perform for privileged operations. This document briefly discusses the organization of kernel calls in MINIX 3 and provides an overview of all kernel calls.

## Organization of kernel calls

A kernel call means that a request is sent to a kernel where it is handled by one of the kernel tasks. The details of assembling a request message, sending it to the kernel, and awaiting the response are conveniently hidden in a system library. The header file of this library is src/include/minix/syslib.hand its implementation is found in src/lib/syslib.

The actual implementation of the kernel calls is defined in one of the kernel tasks. In contrast to MINIX 2, the CLOCK task no longer accepts system calls. Instead, all calls are now directed to the SYSTEM task. Suppose that a program makes a **sys\_call()** system call. By convention, this call is transformed into a request message with type **SYS\_CALL** that is sent to the kernel task SYSTEM. The SYSTEM task handles the request in a function named **do\_call()** and returns the result.

The mapping of kernel call numbers and handler functions is done during the SYSTEM task's intialization in *src/kernel/system.c.* The prototypes of the handler functions are declared in *src/kernel/system.h.* Their implementation is contained in separate files in the directory *src/kernel/system/.* These files are compiled into a library *src/kernel/system/system.a* that is linked with the kernel.

The kernel call numbers and their request and response parameters are defined in *src/include/minix/com.h.* Unfortunately, MINIX 2 does not follow a strict naming scheme. Therefore, numerous message types and parameters have been renamed in MINIX 3. Kernel calls now all start with SYS\_ and all parameters that belong to the same kernel call now share a common prefix.

# Overview of kernel calls in MINIX 3

A concise overview of the kernel calls in MINIX 3 is given in Figure 1. The status of each call compared to MINIX 2 is given in the last column.

Kernel call	Purpose	Status
	PROCESS MANAGEMENT	
SYS_FORK	Fork a process; copy parent process	
SYS_EXEC	Execute a process; initialize registers	
SYS_EXIT	Exit a user process; clear process slot	U
SYS_NICE	Change priority of a user process	Ν
SYS_PRIVCTL	Change system process privileges	Ν
SYS_TRACE	Trace or control process execution	
	SIGNAL HANDLING	
SYS_KILL	Send a signal to a process	U
SYS_GETKSIG	Check for pending kernel signals	Ν
SYS_ENDKSIG	Tell kernel signal has been processed	
SYS_SIGSEND	Start POSIX-style signal handler	
SYS_SIGRETURN	Return from POSIX-style signal	
	MEMORY MANAGEMENT	
SYS_NEWMAP	Install new or updated memory map	
SYS_SEGCTL	Add extra, remote memory segment	Ν
SYS_MEMSET	Write a pattern into physical memory area	Ν
	COPYING DATA	
SYS_UMAP	Map virtual to physical address	U
SYS_VIRCOPY	Copy data using virtual addressing	U
SYS_PHYSCOPY	Copy data using physical addressing	U
SYS_VIRVCOPY	Handle vector with virtual copy requests	U
SYS_PHYSVCOPY	Handle vector with physical copy requests	Ν
	DEVICE I/O	
SYS_DEVIO	Read or write a single device register	Ν
SYS_SDEVIO	Input or output an entire data buffer	Ν
SYS_VDEVIO	Process vector with multiple requests	Ν
SYS_IRQCTL	Set or reset an interrupt policy	Ν
SYS_INT86	Make a real-mode BIOS call	Ν
SYS_IOPENABLE	Give process I/O privilege	Ν
	SYSTEM CONTROL	
SYS_ABORT	Abort MINIX: shutdown the system	U
SYS_GETINFO	Get a copy system info or kernel data	Ν
	CLOCK FUNCTIONALITY	
SYS_SETALARM	Set or reset a synchronous alarm timer	U
SYS_TIMES	Get process times and uptime since boot	U

Figure 1: This figure provides an overview of the kernel calls in MINIX 3. The legenda for the status is: to be <u>New or Updated</u> (i.e., fully revised—all calls got minor updates) since MINIX 2.

## MINIX 3 kernel call interface

MINIX' kernel call interface is detailed below. For each kernel call the message type, the purpose, message type, request and/ or response parameters, and return value are specified. The shorthand Additional remarks about the future status of the call also may be provided.

## Legenda

CONSTANT: defined constant; a number indicating the request type or status *PARAMETER*: message parameter; a field in the request or response message void sys\_call(arguments): system library function; shorthand to make a kernel call

## Alphabetical overview

SYS\_ABORT: Shutdown MINIX and return to the boot monitor—if possible. This is used by PM, FS and TTY. Normal aborts usually are initiated by the user, for example, by means of the 'shutdown' command or typing a 'Ctrl-Alt-Del'. MINIX will also be taken down if a fatal error occurs in the PM or FS.

## $request\ parameters$

ABRT\_HOW: How to abort. One of the values defined in src/include/unistd.h:

- **RBT\_HALT** Halt MINIX and return to the boot monitor.
- **RBT\_REBOOT** Reboot MINIX.
- **RBT\_PANIC** A kernel panic occurred.
- **RBT\_MONITOR** Run the specified code at the boot monitor.
- **RBT\_RESET** Hard reset the system.

*ABRT\_MON\_PROC*: Process to get the boot monitor parameters from. *ABRT\_MON\_LEN*: Length of the boot monitor parameters. *ABRT\_MON\_ADDR*: Virtual address of the parameters.

return value

**OK**: The shutdown sequence was started.

ENIVAL: Invalid process number.

EFAULT: Illegal monitor parameters address.

E2BIG: Monitor parameters exceed maximum length.

## library functions

int sys\_abort(int shutdown\_status, ...);

SYS\_DEVIO: Perform device I/O on behalf of a user-space device driver. The driver can request a single port to be read or written with this call. Also see the SYS\_SDEVIO and SYS\_VDEVIO kernel calls.

#### request parameters

 $DIO\_REQUEST:$  Input or output.

• DIO\_INPUT Read a value from DIO\_PORT.

• DIO\_OUTPUT Write DIO\_VALUE to DIO\_PORT.

 $DIO\_TYPE:$  A flag indicating the type of values.

- DIO\_BYTE Byte type.
- DIO\_WORD Word type.
- DIO\_LONG Long type.

*DIO\_PORT*: The port to be read or written.

DIO\_VALUE: Value to write to the given port. For DIO\_OUTPUT only.

### response parameters

DIO\_VALUE: Value that was read from the given port. For DIO\_INPUT only.

## $return \ value$

 $\mathsf{OK}:$  The port I/O was successfully done.

 $\mathsf{EINVAL}:$  An invalid  $\mathsf{DIO\_REQUEST}$  or  $\mathsf{DIO\_TYPE}$  was provided.

## library functions

int sys\_in(port\_t port, unsigned long value, int io\_type); int sys\_inb(port\_t port, u8\_t \*byte); int sys\_inw(port\_t port, u16\_t \*word); int sys\_inl(port\_t port, u32\_t \*long); int sys\_out(port\_t port, unsigned long \*value, int io\_type); int sys\_outb(port\_t port, u8\_t byte); int sys\_outw(port\_t port, u16\_t word); int sys\_outl(port\_t port, u32\_t long);

SYS\_ENDKSIG: Finish a kernel signal. The PM uses this call to indicate it has processed the kernel signals in the map obtained through a SYS\_GETKSIG kernel call.

### $response\ parameters$

SIG\_PROC: The process that it concerns.

## $return \ value$

 $\mathsf{EINVAL}:$  The process had no pending signals or already exited.

OK: The kernel cleared all pending signals.

 $library\ functions$ 

int sys\_endksig(int proc\_nr);

**SYS\_EXEC**: Update a process' registers after a successfull **exec()** POSIX-call. After the FS has copied the binary image into memory, the PM informs the kernel about the new register details.

## $request\ parameters$

*PR\_PROC\_NR*: Process that executed a program.

*PR\_STACK\_PTR*: New stack pointer.

*PR\_IP\_PTR*: New program counter.

PR\_NAME\_PTR: Pointer to name of program.

## $return\ value$

OK: This call always succeeds.

library functions

int sys\_exec(int proc, char \*stack\_ptr, char \*prog\_name, vir\_bytes pc);

**SYS\_EXIT**: Clear a process slot. This is usually called by the PM to clean up after a user process exited. System processes, including the PM, can also directly call this function to exit themselves.

## request parameters

*PR\_PROC\_NR*: Slot number of exiting process if caller is PM. Use **SELF** to exit the PM.

## $return \ value$

OK: The cleanup succeeded.

EINVAL: Incorrect process number.

EDONTREPLY: This call does not return if a system process exited.

#### library functions

int sys\_exit(int proc\_nr);

SYS\_FORK: Allocate a new (child) process in the kernel process table and initialize it based on the prototype (parent) process. The PM has found a free process slot for the child process in its own process table and now requests the kernel to update the kernel's process table.

## $request\ parameters$

*PR\_PROC\_NR*: Child's process table slot.

*PR\_PPROC\_NR*: Parent, the process that forked.

## $return\ value$

OK: A new process slot was successfully assigned.

EINVAL: Invalid parent process number or child slot in use.

## library functions

int sys\_fork(int parent\_proc\_nr, int child\_proc\_nr);

SYS\_GETINFO: Obtain a copy of a kernel data structure. This call supports user-space device drivers and servers that need certain system information.

## $request\ parameters$

*LREQUEST*: The type of system information that is requested.

- GET\_IMAGE Copy boot image table.
- GET\_IRQHOOKS Copy table with interrupt hooks.
- GET\_KINFO Copy kernel information structure.
- GET\_KMESSAGES Copy buffer with diagnostic kernel messages.
- GET\_LOCKTIMING Copy lock times—if DEBUG\_TIME\_LOCKS is set.
- GET\_MACHINE Copy system environment.
- GET\_MONPARAMS Copy parameters set by the boot monitor.
- GET\_PRIVTAB Copy system privileges table.
- GET\_PROCTAB Copy entire kernel process table.
- GET\_PROC Copy single process table slot.
- GET\_RANDOMNESS Copy randomness gathered by kernel events.

• GET\_SCHEDINFO Copy ready queues and process table.

 $I\_VAL\_PTR:$  Virtual address where the information should be copied to.

 $I\_VAL\_LEN:$  Maximum length that the caller can handle.

*L\_VAL\_PTR2*: Optional, second address. Used when copying scheduling data. *L\_VAL\_LEN2*: Optional, second length. Overloaded for process number.

 $return \ value$ 

OK: The information request succeeded.

EFAULT: An illegal memory address was detected.

E2BIG: Requested data exceeds the maximum provided by the caller.

library functions

int sys\_getinfo(int request, void \*ptr, int len, void \*ptr2, int len2);

int sys\_getirqhooks(struct irq\_hook \*ptr);

- int sys\_getimage(struct boot\_image \*ptr);
- int sys\_getkinfo(struct kinfo \*ptr);
- int sys\_getkmessages(struct kmessages \*ptr);

int sys\_getlocktimings(struct lock\_timingdata \*ptr);

int sys\_getmachine(struct machine \*ptr);

int sys\_getmonparams(char \*ptr, int max\_len);

int sys\_getprivtab(struct priv \*ptr);

int sys\_getproctab(struct proc \*ptr);

int sys\_getproc(struct proc \*ptr, int proc\_nr);

int sys\_getrandomness(struct randomness \*ptr);

int sys\_getschedinfo(struct proc\* ptr, struct proc \*ptr2);

SYS\_GETKSIG: Checks whether there is a process that must be signaled. This is repeatedly done by the PM after receiving a notification that there are kernel signals pending.

response parameters

SIG\_PROC: Return next process with pending signals or NONE.

SIG\_MAP: Bit map with pending kernel signals.

 $return\ value$ 

OK: This call always succeeds.

library functions

int sys\_getksig(int \*proc\_nr, sigset\_t \*sig\_map);

SYS\_INT86: Make a real-mode BIOS on behalf of a user-space device driver. This temporarily switches from 32-bit protected mode to 16-bit real-mode to access the BIOS calls. It is here to support the BIOS\_WINI device driver.

## request parameters

INT86\_REG86: Address of request at the caller.

 $return \ value$ 

OK: BIOS call successfully done. EFAULT: Invalid request address.

library functions

**SYS\_IOPENABLE**: Enable the CPU's I/O privilege level bits for to the given process, so that is allowed to directly perform I/O in user space.

request parameters

*PROC\_NR*: The process to give I/O privileges.

 $return \ value$ 

OK: Always succeeds.

 $library\ functions$ 

SYS\_IRQCTL: Set or reset a hardware interrupt policy for a given IRQ line and enable or disable interrupts for this line. This call allows user-space device drivers to grab a hook for use with the kernel's generic interrupt handler. The kernel's interrupt handler merely notifies the driver about the interrupt with a HARD\_INT message and reenables the IRQ line if the policy says so. The notification message will contain the 'id' provided by the caller as an argument. Once a policy is in place, drivers can enable and disable interrupts.

## $request\ parameters$

 $IRQ\_REQUEST:$  Interrupt control request to perform.

- IRQ\_SETPOLICY Set interrupt policy for the generic interrupt handler.
- IRQ\_RMPOLICY Remove a previously set interrupt policy.
- IRQ\_ENABLE Enable IRQs for the given IRQ line.
- IRQ\_DISABLE Disable IRQs for the given IRQ line.

 $IRQ\_VECTOR:$  IRQ line that must be controlled.

*IRQ\_POLICY*: Bit map with flags indicating IRQ policy.

*IRQ\_HOOK\_ID*: When setting a policy this provides index sent to caller upon interrupt. For other requests it is the kernel hook identifier returned by the kernel.

### response parameters

*IRQ\_HOOK\_ID*: Kernel hook identifier associated with the driver.

return value

EINVAL: Invalid request, IRQ line, hook id, or process number. EPERM: Only owner of hook can toggle interrupts or release the hook.

 $\ensuremath{\mathsf{ENOSPC}}$  : No free IRQ hook could be found.

**OK**: The request was successfully handled.

## library functions

int sys\_irqctl(int request, int irq\_vec, int policy, int \*hook\_id);

int sys\_irqsetpolicy(int irq\_vec, int policy, int \*hook\_id);

int sys\_irqrmpolicy(int irq\_vec, int \*hook\_id);

int sys\_irqenable(int hook\_id);

int sys\_irqdisable(int hook\_id);

SYS\_KILL: Signal a process on behalf of a system server. A system process can signal another process with this call. The kernel notifies the PM about the pending signal for further processing. (Note that the kill() POSIX-call is directly handled at the PM.) The PM uses this call to indirectly send a signal message to a system process. This happens when a signal arrives for a system process that set the special SIG\_MESS signal handler with the sigaction() POSIX-call.

request parameters

SIG\_PROC\_NR: Process to be signaled.

 $\it SIG\_NUMBER:$  Signal number. Range from 0 to  $\_NSIG.$ 

## $return\ value$

 $\mathsf{OK}:$  Call succeeded.

 $\mathsf{EINVAL}:$  Illegal process or signal number.

**EPERM**: Cannot send a signal to a kernel task. PM cannot signal a user process with a notification message.

library functions

int sys\_kill(int proc\_nr, int sig\_nr);

SYS\_MEMSET: Write a 4-byte pattern into the indicated memory area. The call is used by the PM to zero the BSS segment on an exec() POSIX-call. The kernel is ask to do the work for performance reasons.

## $request\ parameters$

MEM\_PTR: Physical base address of the memory area.MEM\_COUNT: Length in bytes of the memory area.MEM\_PATTERN: The 4-byte pattern to be written.

return value

OK: Call always succeeds.

library functions

int sys\_memset(long pattern, phys\_bytes base, phys\_bytes length);

**SYS\_NEWMAP**: Install a new memory map for a newly forked process or if a process' memory map is changed. The kernel fetches the new memory map from PM and updates its data structures.

#### $request\ parameters$

 $PR\_PROC\_NR:$  Install new map for this process.

 $PR\_MEM\_PTR:$  Pointer to memory map at PM.

## $return \ value$

OK: New map was successfully installed. EFAULT: Incorrect address for new memory map.

 $\mathsf{EINVAL}:$  Invalid process number.

## library functions

int sys\_newmap(int proc\_nr, struct mem\_map \*ptr);

SYS\_NICE: Change a process' priority. This is done by passing a nice values between PRIO\_MIN (negative) and PRIO\_MAX (positive). A nice value of zero resets the priority to the default.

#### request parameters

*PR\_PROC\_NR*: Process who's priority should be changed *PR\_PRIORITY*: New nice value for process' priority

 $return\ value$ 

OK: New priority was successfully set.EINVAL: Invalid process number or priority.EPERM: Cannot change priority of kernel task.

El El Milli Camillot change priority

library functions

int sys\_nice(int proc\_nr, int priority);

SYS\_PHYSCOPY: Copy data using physical addressing. The source and/ or destination address can be virtual like with SYS\_VIRCOPY, but in addition an arbitrary physical address is accepted with PHYS\_SEG.

 $request\ parameters$ 

CP\_SRC\_SPACE: Source segment.
CP\_SRC\_ADDR: Virtual source address
CP\_SRC\_PROC\_NR: Process number of the source process.
CP\_DST\_SPACE: Destination segment.
CP\_DST\_ADDR: Virtual destination address
CP\_DST\_PROC\_NR: Process number of the destination process.
CP\_NR\_BYTES: Number of bytes to copy.

#### $return \ value$

OK: The copying was done. EDOM: Invalid copy count.

EFAULT: Virtual to physical mapping failed.

EINVAL: Incorrect segment type or process number.

EPERM: Only owner of REMOTE\_SEG can copy to or from it.

#### library functions

int sys\_abscopy(phys\_bytes src\_phys, phys\_bytes dst\_phys, phys\_bytes count); int sys\_physcopy(int src\_proc, int src\_seg, vir\_bytes src\_vir, int dst\_proc, int dst\_seg, vir\_bytes dst\_vir, phys\_bytes count);

SYS\_PHYSVCOPY: Copy multiple block of data using physical addressing. The request vector is fetched from the caller, and each element is handled like a regular SYS\_PHYSCOPY request. Copying continues until all elements have been processed or an error occurs.

## $request\ parameters$

*VCP\_VEC\_SIZE*: Number of elements in request vector. *VCP\_VEC\_ADDR*: Virtual address of request vector at caller.

#### response parameters

VCP\_NR\_OK: Number of elements successfully copied.

## $return \ value$

 $\mathsf{OK}:$  The copying was done.

EDOM: Invalid copy count.

EFAULT: Virtual to physical mapping failed.

EINVAL: Copy vector too large, incorrect segment or invalid process.

EPERM: Only owner of REMOTE\_SEG can copy to or from it.

library functions

int sys\_physvcopy(phys\_cp\_req \*copy\_vec, int vec\_size, int \*nr\_ok);

SYS\_PRIVCTL: Get a private privilege structure and update a process' privileges. This is used to dynamically start system services.

#### request parameters

CTL\_PROC\_NR: Process who's privileges should be updated.

## $return \ value$

**OK**: The calls succeeded.

EINVAL: Invalid process number.

**ENOSPC**: No free privilege structure found.

## remarks

This system call will be extended to provide both better support and security checks for servers or device drivers that must be dynamically loaded. This is future work.

library functions

SYS\_SDEVIO: Perform device I/O on behalf of a user-space device driver. Note that this call supports only byte and word granularity. The driver can request input or output of an entire buffer. Also see the SYS\_DEVIO and SYS\_VDEVIO kernel calls.

### $request\ parameters$

 $DIO\_REQUEST:$  Input or output.

- DIO\_INPUT Read a value from DIO\_PORT.
- DIO\_OUTPUT Write DIO\_VALUE to DIO\_PORT.

*DIO\_TYPE*: A flag indicating the type of values.

- DIO\_BYTE Byte type.
- DIO\_WORD Word type.

DIO\_PORT: The port to be read or written. DIO\_PROC\_NR: Process number where buffer is.

*DIO\_VEC\_ADDR*: Virtual address of buffer.

DIO\_VEC\_SIZE: Number of elements to input or output.

## $response\ parameters$

DIO\_VALUE: Value that was read from the given port. For DIO\_INPUT only.

#### return value

OK: The port I/O was successfully done. EINVAL: Invalid request or port granularity. EPERM: Cannot do I/O for kernel tasks. EFAULT: Invalid virtual address of buffer.

### library functions

int sys\_insb(port\_t port, u8\_t buffer, int count);

int sys\_insw(port\_t port, u16\_t buffer, int count);

int sys\_outsb(port\_t port, u8\_t buffer, int count);

int sys\_outsw(port\_t port, u16\_t buffer, int count);

int sys\_sdevio(int req, long port, int io\_type, void \*buffer, int count);

SYS\_SEGCTL: Add a memory segment to the process' LDT and its remote memory map. The call returns a selector and offset that can be used to directly reach the remote segment, as well as an index into the remote memory map that can be used with the SYS\_VIRCOPY kernel call.

## $request\ parameters$

SEG\_PHYS: Physical base address of segment. SEG\_SIZE: Size of segment.

#### response parameters

SEG\_INDEX: Index into remote memory map. SEG\_SELECT: Segment selector for LDT entry. SEG\_OFFSET: Offset within segment. Zero, unless 4K granularity is used.

#### return value

ENOSPC: No free slot in remote memory map and LDT.

OK: Segment descriptor successfully added.

## library functions

int sys\_segctl(int \*index, u16\_t \*seg, vir\_bytes \*off, phys\_bytes phys, vir\_bytes size);

**SYS\_SIGRETURN**: Return from a POSIX-style signal handler. The PM requests the kernel to put things in order before the signalled process can resume execution. Also see the **SYS\_SIGSEND** kernel call that pushes a signal context frame onto the stack.

## request parameters

SIG\_PROC: Indicates the process that was signaled.

SIG\_CTXT\_PTR: Pointer to context structure for POSIX-style signal handling.

#### $response\ parameters$

SIG\_PROC: Return next process with pending signals or NONE.

#### $return \ value$

OK: Signal handling action successfully performed.

 $\mathsf{EINVAL}:$  Invalid process number or context structure.

EFAULT: Invalid context structure address, or could not copy signal frame.

#### library functions

int sys\_sigreturn(int proc\_nr, struct sigmsg \*sig\_context);

SYS\_SIGSEND: Signal a process on behalf of the PM by placing the context structure onto the stack. The kernel fetches the structure, initializes it, and copies it to the user's stack.

## $request\ parameters$

 $SIG\_PROC:$  Indicates the process that was signaled.

SIG\_CTXT\_PTR: Pointer to context structure for POSIX-style signal handling.

response parameters

SIG\_PROC: Return next process with pending signals or NONE.

return value

OK: Signal handling action successfully performed.

EINVAL: Invalid process number.

EPERM: Cannot signal kernel tasks.

EFAULT: Invalid context structure address, or could not copy signal frame.

library functions

int sys\_sigsend(int proc\_nr, struct sigmsg \*sig\_context);

SYS\_SETALARM : Set or reset a synchronous alarm timer. When the timer expires it causes a SYN\_ALARM notification message with the current uptime as an argument to be sent to the caller. Only system processes can request synchronous alarms.

request parameters

*ALRM\_EXP\_TIME*: Absolute or relative expiration time in ticks for this alarm. *ALRM\_ABS\_TIME*: Zero if expire time is relative to the current uptime.

ALIAM\_ADD\_TIME. Leto it expire time is relative to the curr

 $response\ parameters$ 

ALRM\_TIME\_LEFT: Ticks left on the previous alarm.

 $return\ value$ 

OK: The alarm was successfully set.

EPERM: User processes cannot request alarms.

library functions

int sys\_setalarm(clock\_t expire\_time, int abs\_flag);

SYS\_TIMES: Get the kernel's uptime since boot and process execution times.

#### request parameters

*T\_PROC\_NR*: The process to get the time information for, or NONE.

 $response\ parameters$ 

*T\_USER\_TIME*: Process' user time in ticks, if valid number. *T\_SYSTEM\_TIME*: Process' system time in ticks, if valid number. *T\_BOOT\_TICKS*: Number of ticks since MINIX boot.

## $return \ value$

OK: Always succeeds.

## $library\ functions$

int sys\_times(int proc\_nr, clock\_t \*ptr);

SYS\_TRACE: Monitor or control execution of the given process. Handle the debugging commands supported by the ptrace() system call.

## $request\ parameters$

CTL\_REQUEST: The tracing request.

- T\_STOP Stop the process.
- **T\_GETINS** Return value from instruction space.
- T\_GETDATA Return value from data space.
- T\_GETUSER Return value from user process table.

- T\_SETINS Set value from instruction space.
- T\_SETDATA Set value from data space.
- **T\_SETUSER** Set value in user process table.
- T\_RESUME Resume execution.
- T\_STEP Set trace bit.

CTL\_PROC\_NR: The process number that is being traced.

CTL\_ADDRESS: Virtual address in the traced process' space.

*CTL\_DATA*: Data to be written.

## response parameters

*CTL\_DATA*: Data be returned.

 $return \ value$ 

OK: Trace operation succeeded.

EIO: Set or get value failed.

EINVAL: Unsupported trace request.

PERM: Can only trace user processes.

#### library functions

int sys\_trace(int request, int proc\_nr, long addr, long \*data\_ptr);

SYS\_UMAP: Map a virtual address to a physical address and return the physical address. The virtual address can be in LOCAL\_SEG, REMOTE\_SEG, or BIOS\_SEG. An offset in bytes can be passed to verify whether this also falls within the segment.

## $request\ parameters$

CP\_SRC\_PROC\_NR: Process number of the address relates to.

CP\_SRC\_SPACE: Segment identifier.

CP\_SRC\_ADDR: Offset within segment.

CP\_NR\_BYTES: Number of bytes from start.

## $response\ parameters$

CP\_DST\_ADDR: Physical address if mapping succeeded.

## $return \ value$

OK: The copying was done.

EFAULT: Virtual to physical mapping failed.

EINVAL: Incorrect segment type or process number.

## remarks

Address zero within **BIOS\_SEG** returns **EFAULT**, while it the zeroth BIOS interrupt vector in fact is a valid address.

## library functions

int sys\_umap(int proc\_nr, int seg, vir\_bytes vir\_addr, vir\_bytes count, phys\_bytes \*phys\_addr);

SYS\_VDEVIO: Perform a series of device I/O on behalf of a user process. The call accepts a pointer to an array of (port,value)-pairs that is to be handeld at once. Hardware interrupts are temporarily disabled to prevented the bactch of I/O calls to be interrupted. Also see SYS\_DEVIO and SYS\_SDEVIO.

## $request\ parameters$

DIO\_REQUEST: Input or output.

- DIO\_INPUT Read a value from DIO\_PORT.
- DIO\_OUTPUT Write DIO\_VALUE to DIO\_PORT.

DIO\_TYPE: A flag indicating the type of values.

- DIO\_BYTE Byte type.
- DIO\_WORD Word type.
- DIO\_LONG Long type.

*DIO\_VEC\_SIZE*: The number of ports to be handled.

DIO\_VEC\_ADDR: Virtual address of the (port, value)-pairs in the caller's space.

## return value

OK: The port I/O was successfully done.

EINVAL: Invalid request or granularity.

E2BIG: Vector size exceeds maximum that can be handled.

EFAULT: The address of the (port, value)-pairs is erroneous.

## library functions

int sys\_voutb(pvb\_pair\_t \*pvb\_vec, int vec\_size); int sys\_voutw(pvw\_pair\_t \*pvw\_vec, int vec\_size); int sys\_voutl(pvl\_pair\_t \*pvl\_vec, int vec\_size); int sys\_vinb(pvb\_pair\_t \*pvb\_vec, int vec\_size); int sys\_vinw(pvw\_pair\_t \*pvw\_vec, int vec\_size);

int sys\_vinl(pvl\_pair\_t \*pvl\_vec, int vec\_size);

SYS\_VIRCOPY: Copy data using virtual addressing. The virtual can be in three segments: LOCAL\_SEG (text, stack, data segments), REMOTE\_SEG (e.g., RAM disk, video memory), and the BIOS\_SEG (BIOS interrupt vectors, BIOS data area). This is the most common system call relating to copying.

## request parameters

CP\_SRC\_SPACE: Source segment.

CP\_SRC\_ADDR: Virtual source address

CP\_SRC\_PROC\_NR: Process number of the source process.

CP\_DST\_SPACE: Destination segment.

CP\_DST\_ADDR: Virtual destination address

CP\_DST\_PROC\_NR: Process number of the destination process.

CP\_NR\_BYTES: Number of bytes to copy.

## return value

OK: The copying was done.

EDOM: Invalid copy count.

EFAULT: Virtual to physical mapping failed.

EPERM: No permission to use PHYS\_SEG.

EINVAL: Incorrect segment type or process number.

EPERM: Only owner of REMOTE\_SEG can copy to or from it.

#### library functions

int sys\_biosin(vir\_bytes bios\_vir, vir\_bytes dst\_vir, vir\_bytes bytes);

int sys\_biosout(vir\_bytes src\_vir, vir\_bytes bios\_vir, vir\_bytes bytes);

int sys\_datacopy(vir\_bytes src\_proc, vir\_bytes src\_vir, dst\_proc, dst\_vir, vir\_bytes bytes);

int sys\_textcopy(vir\_bytes src\_proc, vir\_bytes src\_vir, dst\_proc, dst\_vir, vir\_bytes bytes);

int sys\_stackcopy(vir\_bytes src\_proc, vir\_bytes src\_vir, dst\_proc, dst\_vir, vir\_bytes bytes); int sys\_vircopy(int src\_proc, int src\_seg, vir\_bytes src\_vir, int dst\_proc, int dst\_seg, vir\_bytes dst\_vir, phys\_bytes bytes);

SYS\_VIRVCOPY: Copy multiple blocks of data using virtual addressing. The request vector is fetched from the caller, and each element is handled like a regular SYS\_VIRCOPY request. Copying continues until all elements have been processed or an error occurs.

 $request\ parameters$ 

VCP\_VEC\_SIZE: Number of elements in request vector.

VCP\_VEC\_ADDR: Virtual address of request vector at caller.

 $response\ parameters$ 

VCP\_VEC\_OK: Number of elements successfully copied.

## $return \ value$

 $\mathsf{OK}:$  The copying was done.

EDOM: Invalid copy count.

EFAULT: Virtual to physical mapping failed.

**EPERM**: No permission to use **PHYS\_SEG**.

EINVAL: Copy vector too large, incorrect segment or invalid process.

EPERM: Only owner of REMOTE\_SEG can copy to or from it.

## library functions

int sys\_virvcopy(vir\_cp\_req \*copy\_vec, int vec\_size, int \*nr\_ok);