Linux 2.6 performance improvement through readahead optimization

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What is the problem?



DSS: Decision Support System

I/O architecture





Results of instrumentation



NOTE: application is generating mostly 256K random reads

Results of the instrumentation cont..

- I/O scheduler could not merge requests because
- It received one-page read requests 50% of the time because...
- READAHEAD did not READ AHEAD!
- Result: Synchronous 1 page reads most of the time.

READAHEAD! What is its job?

- Predict future read requests using past request patterns.
- Asynchronously bring future read requests into the page cache.
- Maximize page-cache hit!
- Minimize synchronous page read.
- Minimize waste of resources.
- Maximize size of the asynchronous read requests.

Readahead algorithm in 2.6.0

Current window



Readahead window



Next window size

MAX_READAHEAD

- Current window tracks pages that satisfy the current read request.
- Readhead window tracks pages that satisfy future read requests.
- Next window size tracks the number of page frames in the next readahead window or the next current window.

Case 1: read request for page 0 size 1



Note: current window is populated

Case 2: read request for page 1



Note: readahead window is populated

Case 3: read request for page 10



Note: next window size increments by 2

Case 4: read request for page 16



Case 5: read request for page 200



- Next window size decremented by 2
- New current window brought in with next widow size page frames
- Readahead window reset

Case 6: read request for page 200



- Reading ahead stopped
- Readahead window reset
- Only the requested page brought in the current window

Case 7: read request for page 201



- Reading ahead stopped
- Only the requested page brought in the current window

Case 8: read request for page 220



- Reading ahead stopped
- current window shrinks down to one page frame :-(

Case 9: read request for page 231



- Reading ahead mode resumed!
- next readahead size set to MAX_READAHEAD!

Case 10: read request for page 210. Pages 232 to 263 are already in the page cache



- Since all the **readahead window** pages already reside in page cache shrink **next readahead size** by 1!
- Populate the readahead window

Problems of 2.6.0 readahead

- Page-0 problem.
- Page-cache-hit on first read problem.
- Lots of wasted pages.

Problem-1 of 2.6.0 readahead

Consider the following file read pattern. 4-Page random read requests.

10, 11, 12, 1390, 91, 92, 9353, 54, 55, 5628, 29, 30, 31

STEP 1: Initial state



STEP 2: Read page 10



STEP 3: Read page 11



STEP 4: Read page 12



STEP 5: Read page 13



STEP 6: Read page 90



STEP 7: Read page 91



STEP 8: Read page 92



STEP 9: Read page 93





STEP 11: Read page 55



STEP 12: Read page 56



STEP 13: Read page 56





STEP 15: Read page 29


STEP 16: Read page 30



• NOTE: synchronous read of page 30

STEP 18: Read page 31



• NOTE: synchronous read of page 31

Where is the bug?

- If file is accessed from offset 0, file access is sequential?
- In the readahead off-mode any miss in contiguity is penalized heavily by resetting the current window.

Problem-2 of 2.6.0 readahead

• If the pages read in the **current-window** the first time, all reside in the page-cache reada-head is turned-off!

Consider the following file read pattern. 4-Page random read requests.

0, 1, 2, 3 90, 91, 92, 93 53, 54, 55, 56 28, 29, 30, 31

STEP 1: Initial state



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STEP 2: Read page 0





Problem-3 of 2.6.0 readahead

 Lots of wasted pages, wasted bandwidth, wasted resources!

Consider the following file read pattern. 4-Page random read requests.

10, 11, 12, 13 90, 91, 92, 93 53, 54, 55, 56 28, 29, 30, 31

STEP 1: Initial state



STEP 1: request for page 10



STEP 2: request for page 11



STEP 3: request for page 12



STEP 4: request for page 13



STEP 5: request for page 90



STEP 6: request for page 91



STEP 7: request for page 92



STEP 8: request for page 93



STEP 9: request for page 53



STEP 10: request for page 54



STEP 11: request for page 55



STEP 12: request for page 56



STEP 13: request for page 28



STEP 14: request for page 29



STEP 15: request for page 30



STEP 16: request for page 31



2.6.7 Readahead algorithm

- No penalty if reads don't start from page 0
- Readahead mode is not switched off for the first read request even if all the pages are found in the page cache.
- The average size of read request is tracked.
- The average size determines the size of the next current window.
- Only if average size is greater than MAX_READAHEAD, readahead window is populated.

2.6.7 readahead

Consider the following file read pattern. 4-Page random read requests.

STEP 1: Initial state



STEP 2: Request for page 10



STEP 3: Request for page 11



STEP 4: Request for page 12



STEP 5: Request for page 13



STEP 6: Request for page 90



STEP 7: Request for page 91



STEP 8: Request for page 92



STEP 9: Request for page 93



STEP 10: Request for page 53


STEP 11: Request for page 54



STEP 12: Request for page 55



STEP 13: Request for page 56



STEP 14: Request for page 28



STEP 15: Request for page 29



STEP 16: Request for page 30



STEP 17: Request for page 31



STEP 18: Request for page 70



STEP 19: Request for page 71



STEP 20: Request for page 72



STEP 21: Request for page 73



2.6.0 v/s 2.6.7 Application generates 16-page request followed by 117-page seek, repeated 100 times. First read request is not at offset 0.

random workload: request-blocksize 16pages, seek between reads 117pages



2.6.0 v/s 2.6.7

Application generates 4-page request followed by 96-page seek, repeated 1000 times.



2.6.0 v/s 2.6.7

Application generates sequential 16-page request repeated 1000 times. [SEQUENTIAL READ PATTERN]

sequential workload: request-blocksize 16pages, seek between reads 0pages



DSS performance comparison



DSS performance comparison



lozone on NFS comparison

	2.6.0 through- put in KB/sec	2.6.7 through- put in KB/sec
Random read	10056.49	19968.79
Random mix read	10053.37	21565.43
Reverse read	10125.13	20138.83
Stride read	7210.96	14461.63
Sequential read	14464.20	13614.49
Sequential re-read	14591.19	13715.94
Pread	11703.76	13668.21

Filesize = 4G, blocksize=128K

Current Limitations

- Too large ramp-up time especially for sequential reads
- For sequential or large-random reads current-window pages and readaheadwindow pages can be brought together
- Slow read path(readahead-off mode) is too slow!
- Multiple threaded reads confuses the algorithm.

Current Limitations cont..

- Readahead algorithm is unaware of the size of the current read even though the File-Mapper knows about this.
- Does not adapt to page-cache pressure.
- MAX_READAHEAD default value is set to 32 which is too small for enterprise workloads.

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2.6.0 readahead cont..

Case 6: read request for page 112



Note: next window size cannot exceed MAX_READAHEAD 94

2.6.0 readahead cont..

Case 5: read request for page 17



Note: readahead window is populated